

# Tractable Query Answering for Expressive Ontologies and Rules\*

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\*Published at ISWC 2017

# Disjunctive Existential Rules

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Facts

# Disjunctive Existential Rules

$\text{HasParent}(x, y) \wedge \text{HasSister}(y, z) \rightarrow \text{HasAunt}(x, z)$

$\text{Human}(x) \rightarrow \exists y . \text{HasParent}(x, y) \wedge \text{Human}(y)$

$\text{Animal}(x) \rightarrow \text{Herbivore}(x) \vee \text{Carnivore}(x) \vee \text{Omnivore}(x)$

$\text{P}(x, a, y) \wedge \text{R}(y, w) \wedge \text{S}(w, x) \rightarrow \exists v . (\text{R}(w, v) \wedge \text{A}(v)) \vee \text{D}(x)$

## Facts

# Disjunctive Existential Rules

$\text{HasParent}(x, y) \wedge \text{HasSister}(y, z) \rightarrow \text{HasAunt}(x, z)$

$\text{Human}(x) \rightarrow \exists y . \text{HasParent}(x, y) \wedge \text{Human}(y)$

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## Facts

$\text{HasFriend}(\text{stan}, \text{kyle})$

$\text{Dead}(\text{kenny})$

$\text{P}(a, c, d)$

# Disjunctive Existential Rules

$\text{HasParent}(x, y) \wedge \text{HasSister}(y, z) \rightarrow \text{HasAunt}(x, z)$

$\text{Human}(x) \rightarrow \exists y . \text{HasParent}(x, y) \wedge \text{Human}(y)$

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$P(x, a, y) \wedge R(y, w) \wedge S(w, x) \rightarrow \exists v . (R(w, v) \wedge A(v)) \vee D(x)$

## Facts

$\text{HasFriend}(\text{stan}, \text{kyle})$

$\text{Dead}(\text{kenny})$

$P(a, c, d)$

*Remark.* If normalised, SROIQ ontologies can be translated into equivalent programs of disjunctive existential rules.

# The Disjunctive Chase

# The Disjunctive Chase

$\text{Director}(x) \rightarrow \exists y . \text{Directs}(x, y) \wedge \text{Film}(y)$

$\text{Film}(x) \rightarrow \exists z . \text{IsDirectedBy}(x, z) \wedge \text{Director}(z)$

$\text{Film}(\text{ai})$



# The Skolem Disjunctive Chase

$\text{Director}(x) \rightarrow \exists y . \text{Directs}(x, y) \wedge \text{Film}(y)$

$\text{Film}(x) \rightarrow \exists z . \text{IsDirectedBy}(x, z) \wedge \text{Director}(z)$

$\text{Film}(\text{ai})$

# The Skolem Disjunctive Chase

$\text{Director}(x) \rightarrow \text{Directs}(x, f_y(x)) \wedge \text{Film}(f_y(x))$

$\text{Film}(x) \rightarrow \text{IsDirectedBy}(x, f_z(x)) \wedge \text{Director}(f_z(x))$

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# The Skolem Disjunctive Chase

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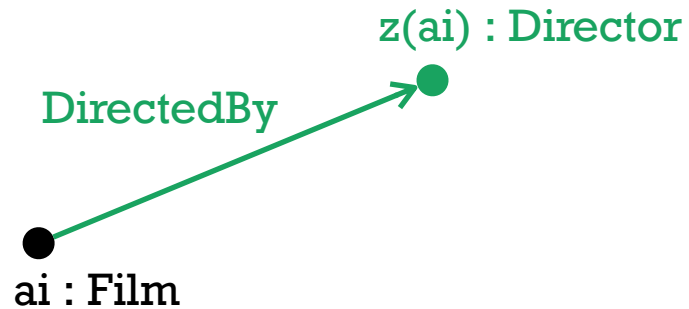
$\text{ai} : \text{Film}$

# The Skolem Disjunctive Chase

$\text{Director}(x) \rightarrow \text{Directs}(x, y(x)) \wedge \text{Film}(y(x))$

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$\text{Film}(a_i)$

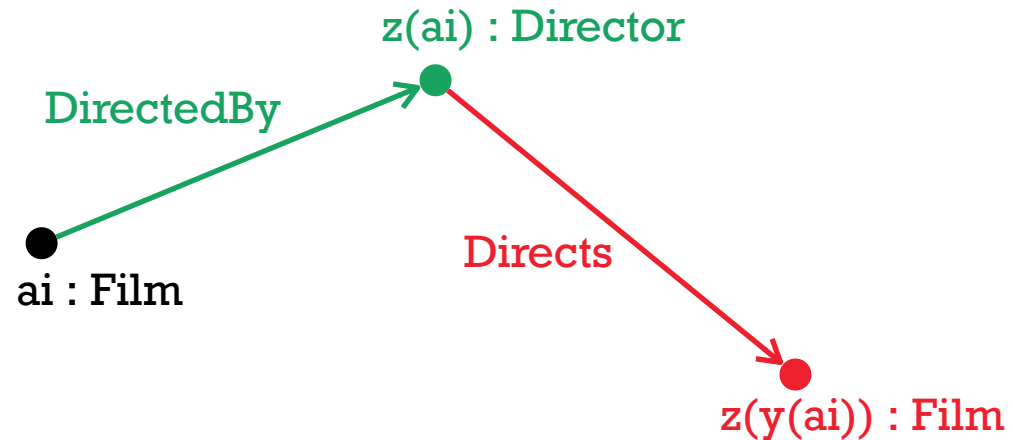


# The Skolem Disjunctive Chase

$\text{Director}(x) \rightarrow \text{Directs}(x, y(x)) \wedge \text{Film}(y(x))$

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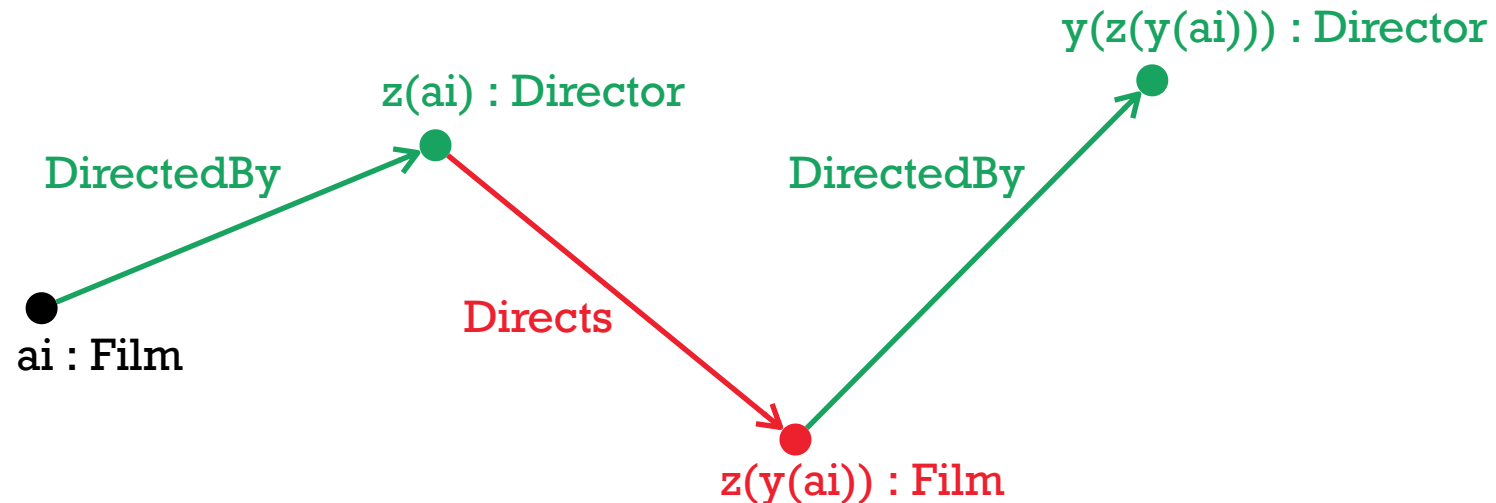


# The Skolem Disjunctive Chase

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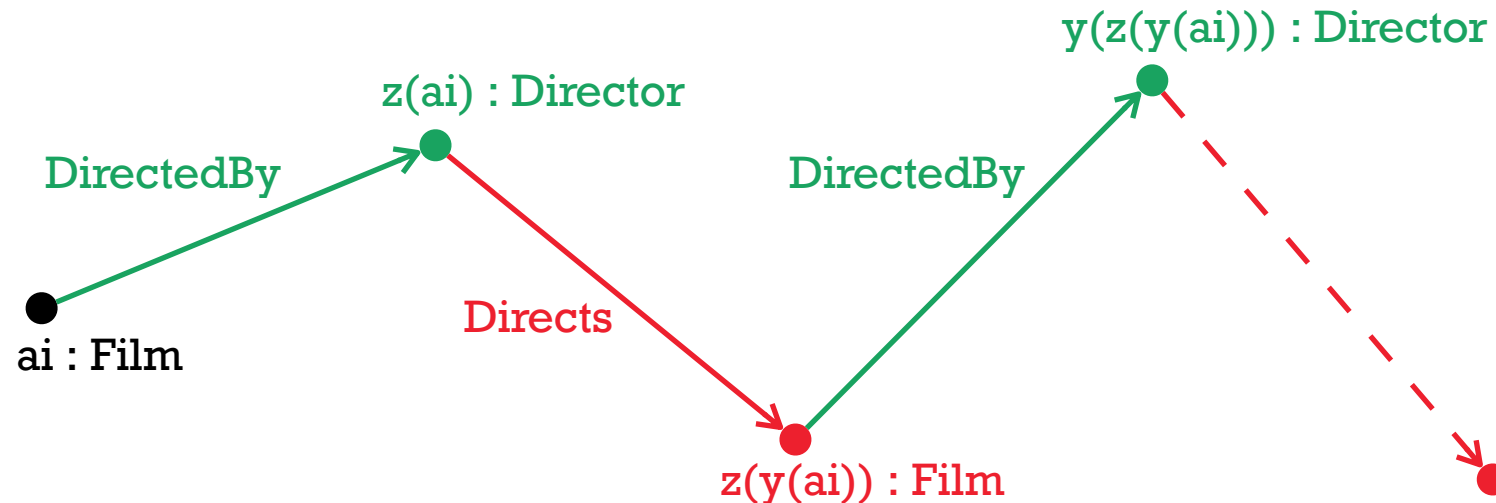


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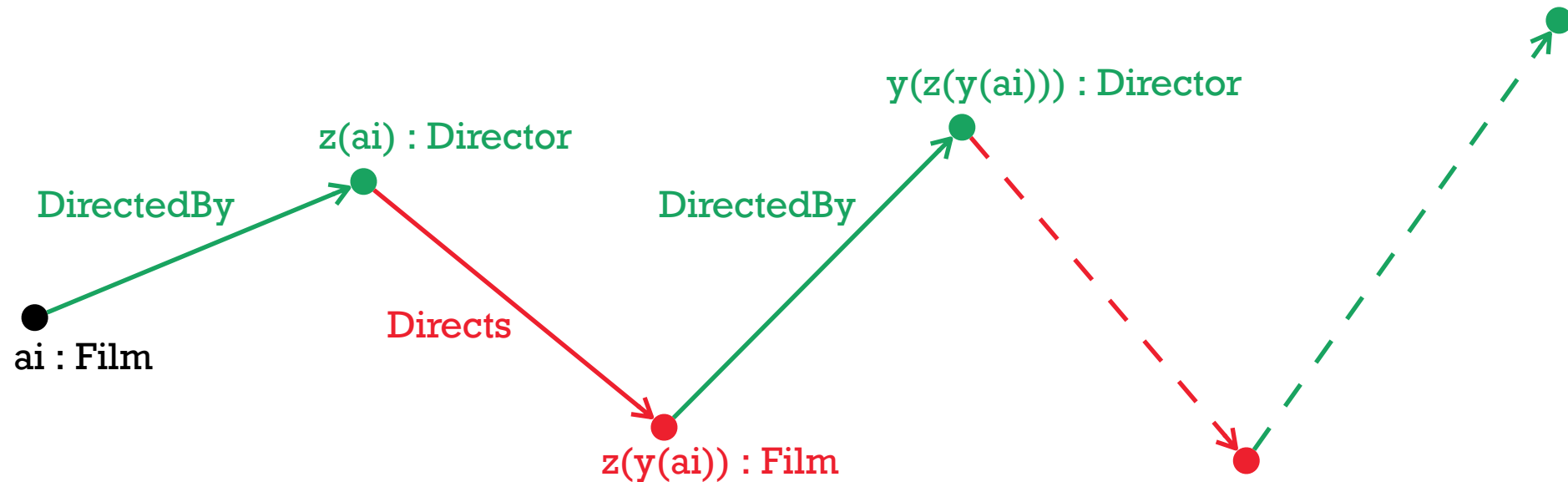


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$\text{Film}(\text{ai})$

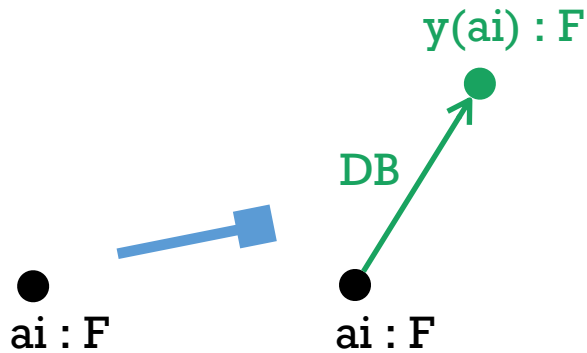
●  
ai : F

# The Skolem Disjunctive Chase

$\text{Director}(x) \rightarrow \text{Directs}(x, y(x)) \wedge \text{Film}(y(x))$

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$\text{Film}(\text{ai})$

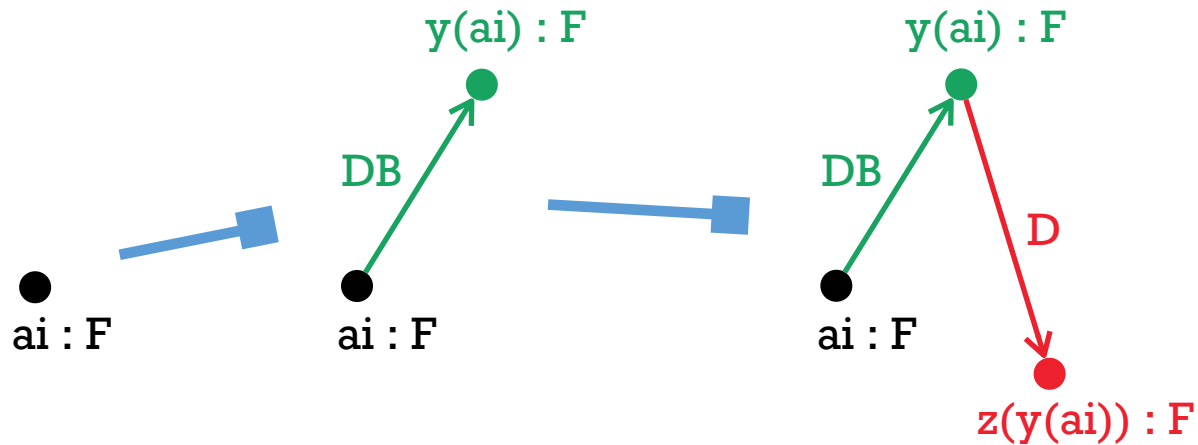


# The Skolem Disjunctive Chase

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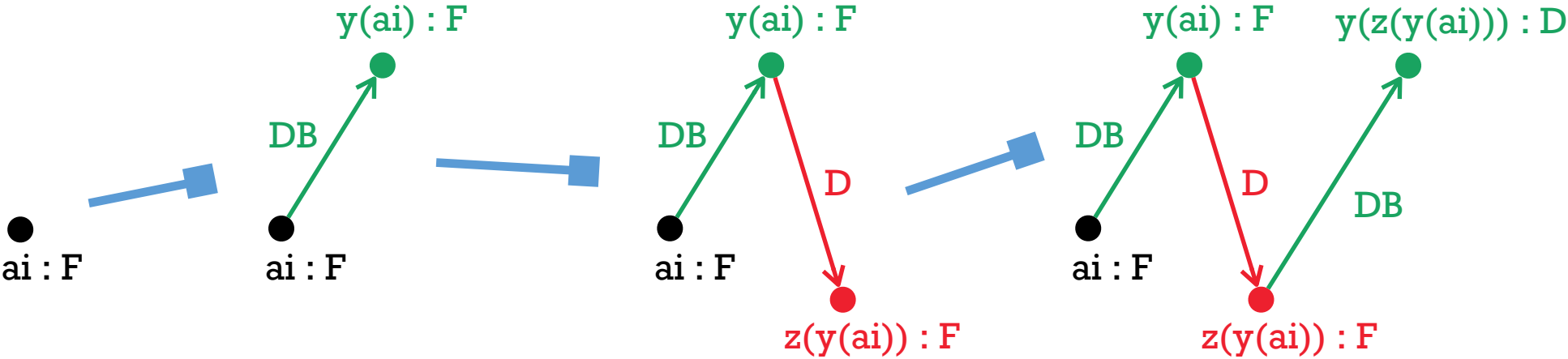


# The Skolem Disjunctive Chase

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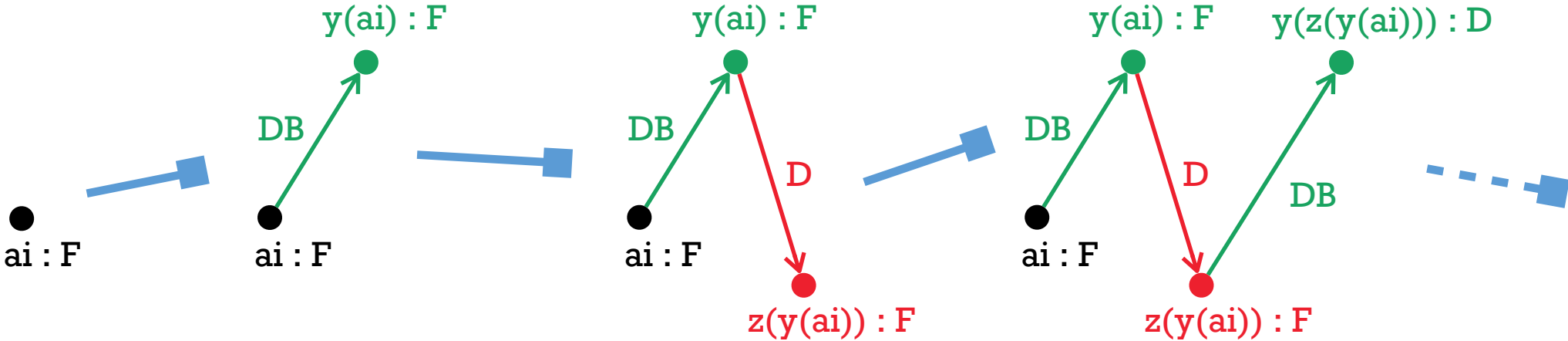


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$\text{Film}(a_i)$



# The Skolem Disjunctive Chase

$\text{Animal}(x) \rightarrow \text{Vertebrate}(x) \vee \text{Invertebrate}(x)$

$\text{Animal}(a) \quad \text{Animal}(b)$



# The Skolem Disjunctive Chase

$\text{Animal}(x) \rightarrow \text{Vertebrate}(x) \vee \text{Invertebrate}(x)$

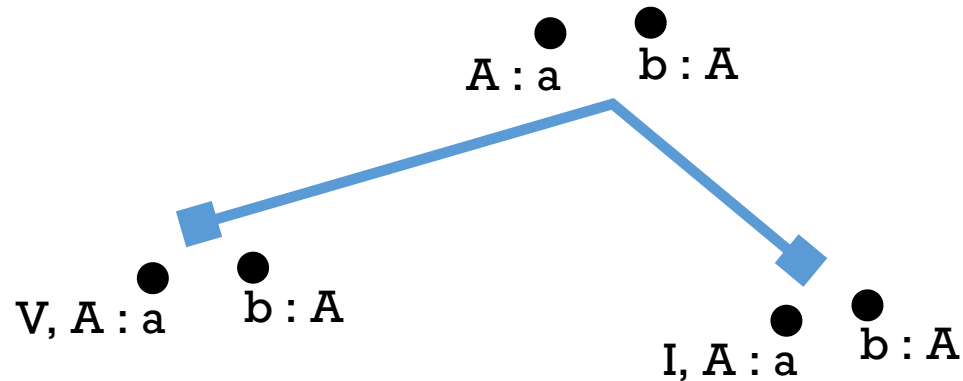
$\text{Animal}(a) \quad \text{Animal}(b)$

$A : a \quad b : A$

# The Skolem Disjunctive Chase

$\text{Animal}(x) \rightarrow \text{Vertebrate}(x) \vee \text{Invertebrate}(x)$

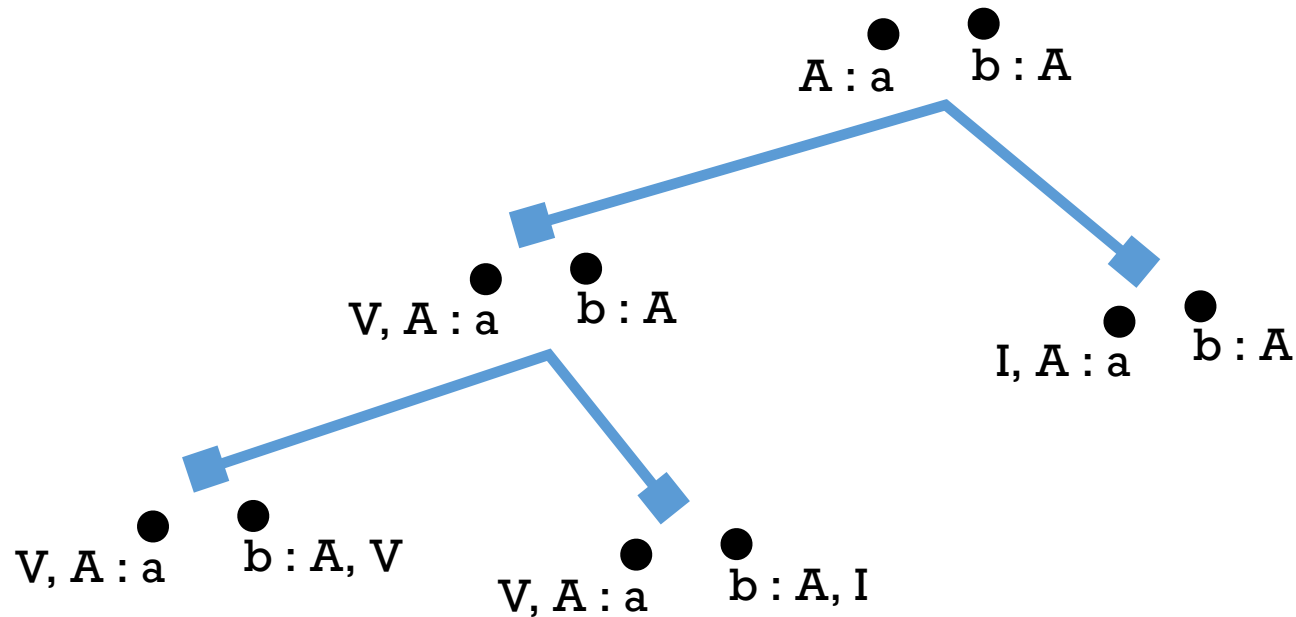
$\text{Animal}(a) \quad \text{Animal}(b)$



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$\text{Animal}(x) \rightarrow \text{Vertebrate}(x) \vee \text{Invertebrate}(x)$

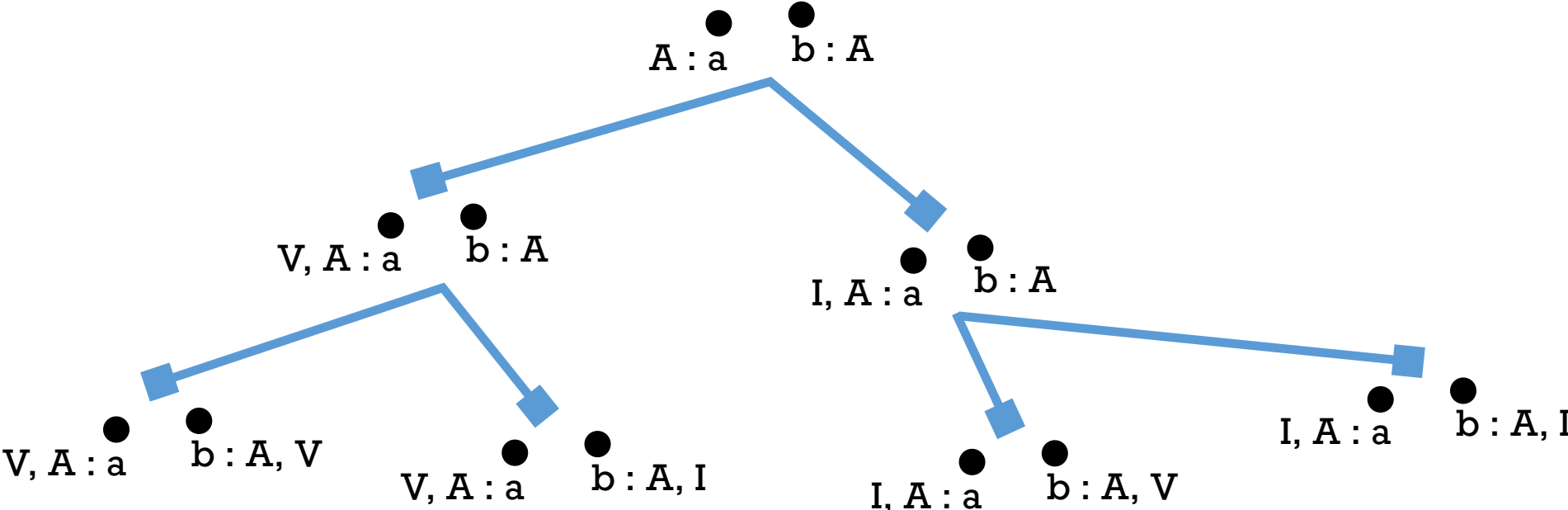
$\text{Animal}(a) \quad \text{Animal}(b)$



# The Skolem Disjunctive Chase

$\text{Animal}(x) \rightarrow \text{Vertebrate}(x) \vee \text{Invertebrate}(x)$

$\text{Animal}(a) \quad \text{Animal}(b)$



# Ensuring Tractability of the Disjunctive Chase

# Existential Dependency Graph

$$A(x) \rightarrow \exists y . S(x, y) \wedge B(y)$$

$$B(x) \rightarrow \exists z . R(x, z) \wedge D(z)$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow \exists w . R(x, w)$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$

# Existential Dependency Graph

$$\mathbf{A}(x) \rightarrow \mathbf{S}(x, y(x)) \wedge \mathbf{B}(y(x))$$

$$\mathbf{B}(x) \rightarrow \mathbf{R}(x, z(x)) \wedge \mathbf{D}(z(x))$$

$$\mathbf{D}(x) \rightarrow \mathbf{E}(x)$$

$$\mathbf{E}(x) \rightarrow \mathbf{R}(x, w(x))$$

$$\mathbf{B}(x) \wedge \mathbf{C}(x) \rightarrow \mathbf{E}(x)$$

$$\mathbf{S}(x, y) \rightarrow \mathbf{C}(x)$$

# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$

w  
●

● z

●  
y



# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$

w  
●

● z

●  
y

# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$

$A(c)$

$S(c, y(c)), B(y(c))$

$R(y(c), z(y(c))), D(z(y(c)))$

w  
●

● z

●  
y

# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

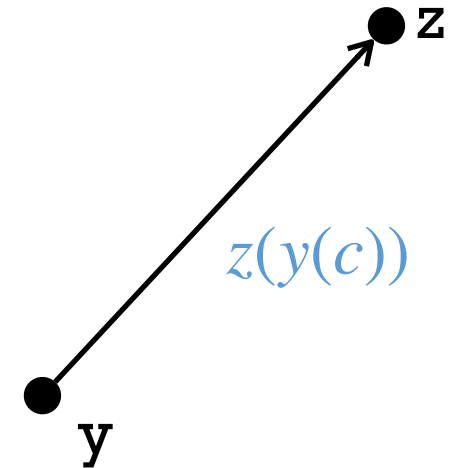
$$S(x, y) \rightarrow C(x)$$

$A(c)$

$S(c, y(c)), B(y(c))$

$R(y(c), z(y(c))), D(z(y(c)))$

$w$   
●



# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

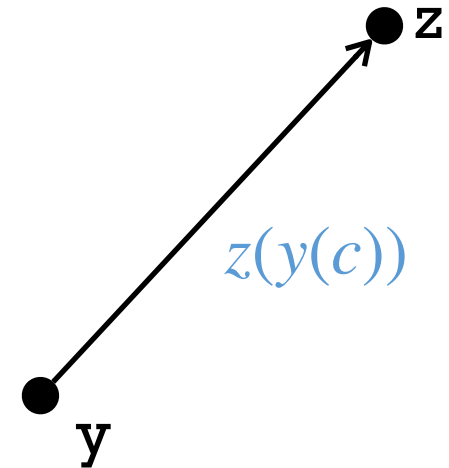
$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$

w  
●



# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

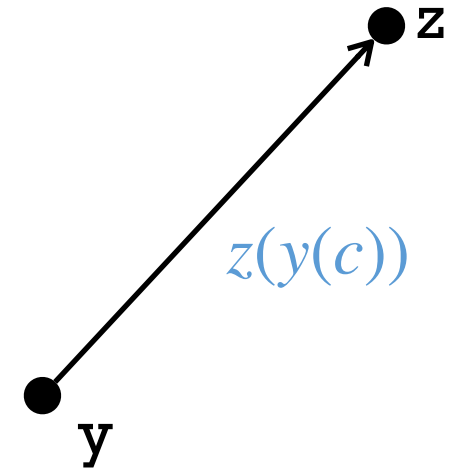
$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$

w  
●



# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$

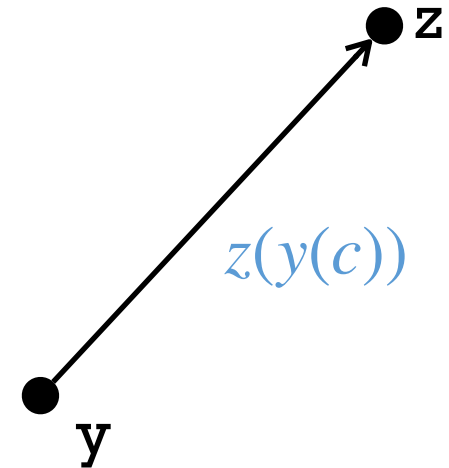
$B(c)$

$R(c, z(c)), D(z(c)),$

$E(z(c)),$

$R(z(c), w(z(c)))$

$w$   
●



# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

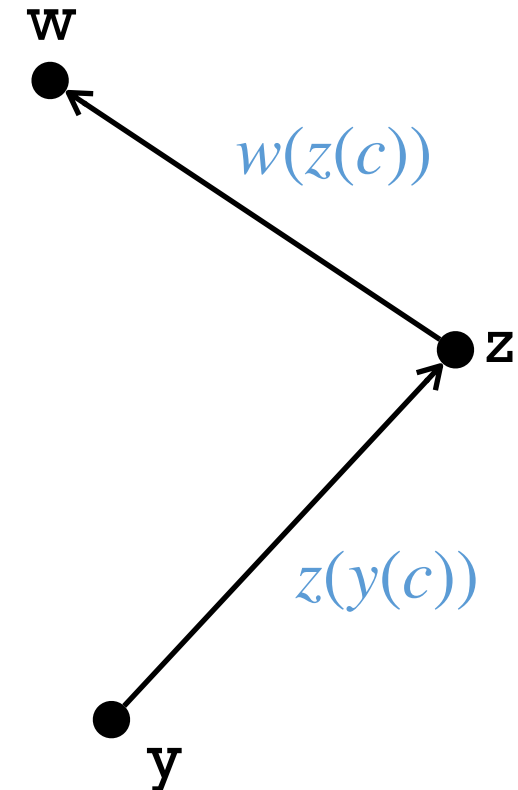
$$S(x, y) \rightarrow C(x)$$

$B(c)$

$R(c, z(c)), D(z(c)),$

$E(z(c)),$

$R(z(c), w(z(c)))$



# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

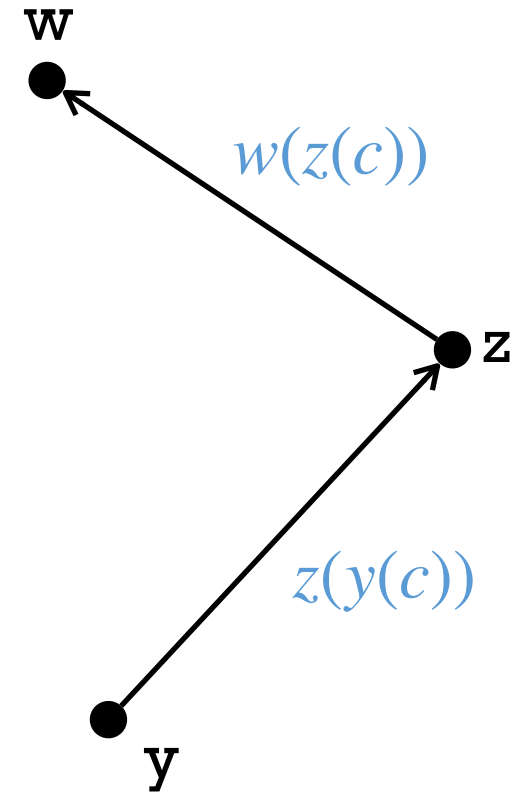
$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$





# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

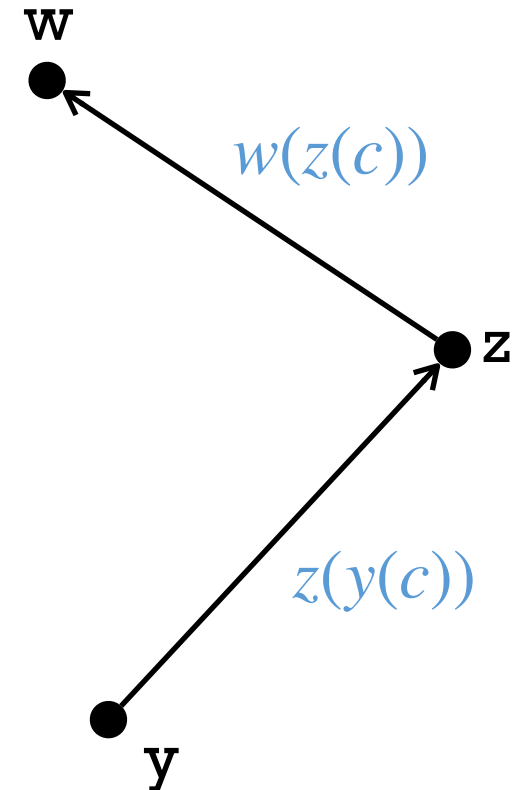
$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$



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$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

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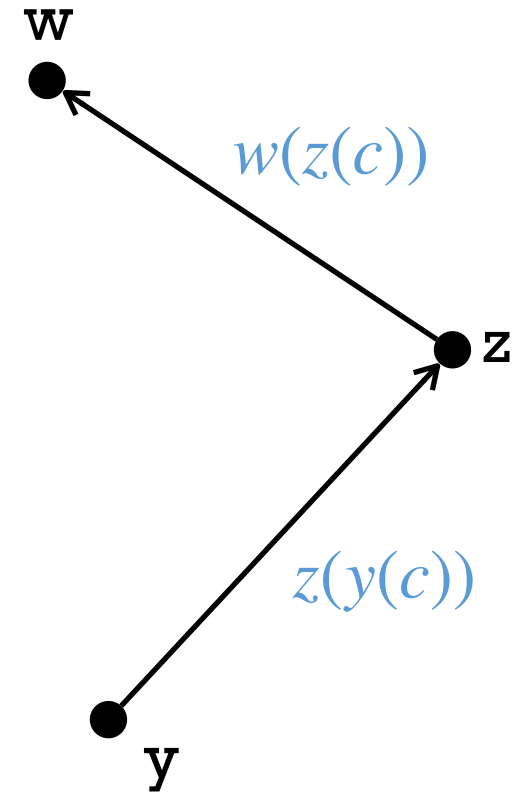
$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$

$A(c)$   
 $S(c, y(c)), B(y(c)),$   
 $C(y(c)),$   
 $E(y(c)),$   
 $R(y(c), w(y(c)))$



# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

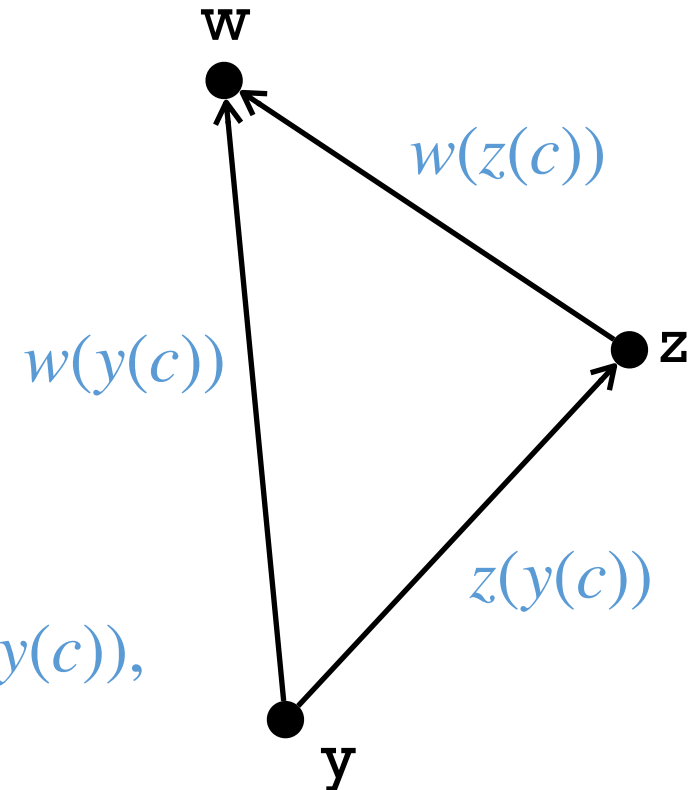
$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$



$A(c)$   
 $S(c, y(c)), B(y(c)),$   
 $C(y(c)),$   
 $E(y(c)),$   
 $R(y(c), w(y(c)))$

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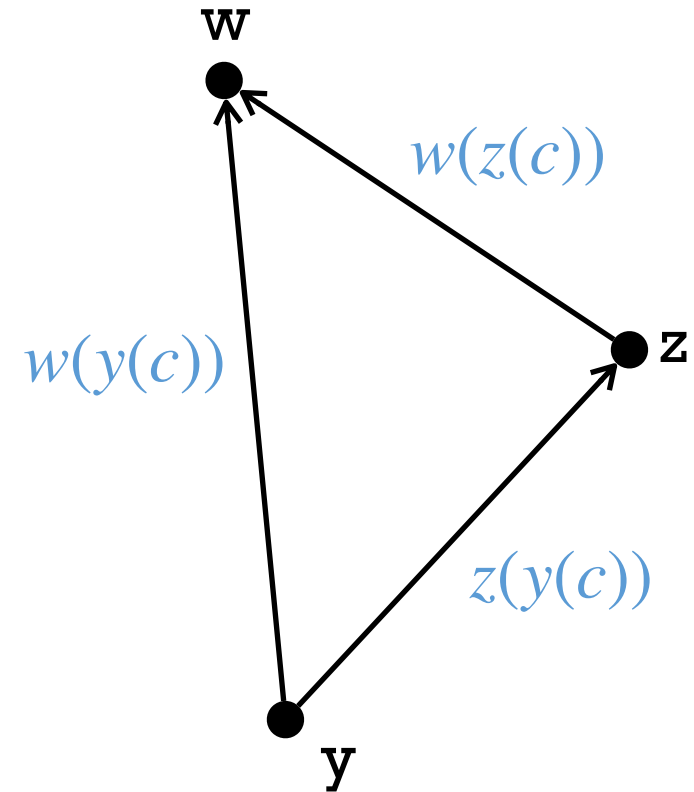
$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

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$$B(x) \wedge C(x) \rightarrow E(x)$$

$$S(x, y) \rightarrow C(x)$$



# Existential Dependency Graph

$$A(x) \rightarrow S(x, y(x)) \wedge B(y(x))$$

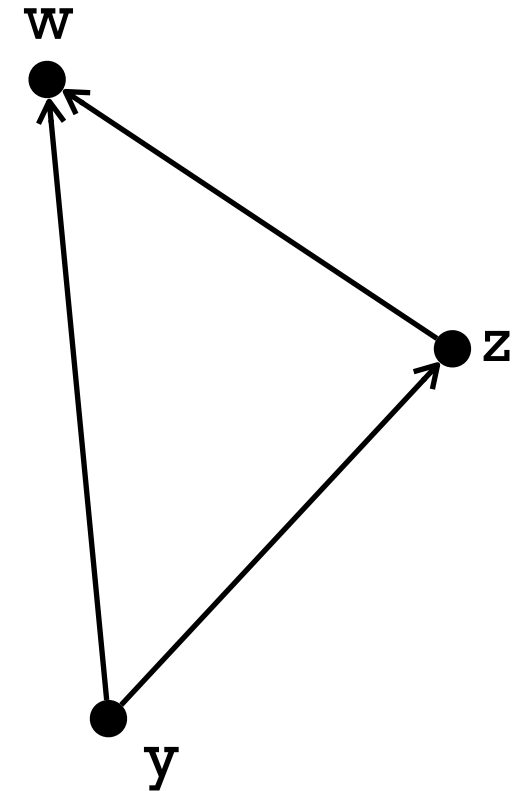
$$B(x) \rightarrow R(x, z(x)) \wedge D(z(x))$$

$$D(x) \rightarrow E(x)$$

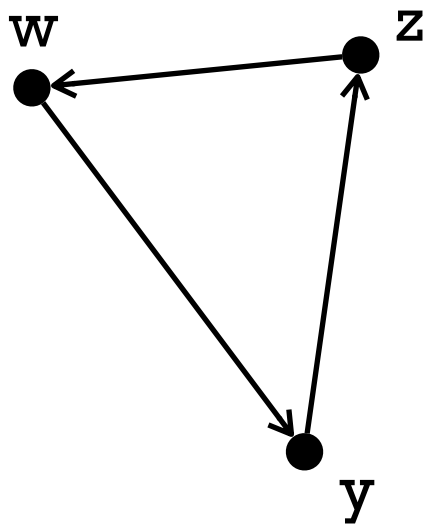
$$E(x) \rightarrow R(x, w(x))$$

$$B(x) \wedge C(x) \rightarrow E(x)$$

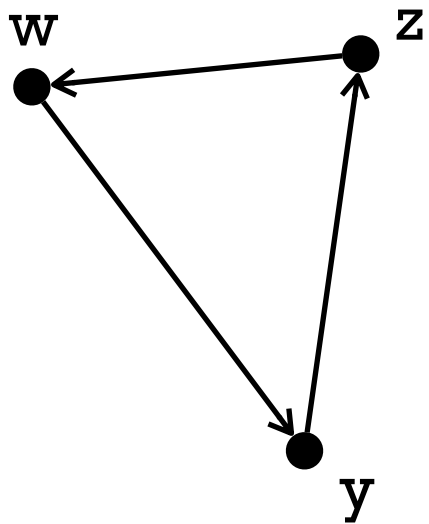
$$S(x, y) \rightarrow C(x)$$



# (a) Acyclicity

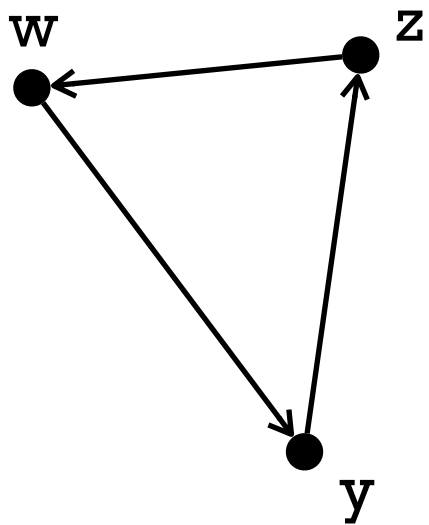


# (a) Acyclicity



$z(c)$

# (a) Acyclicity

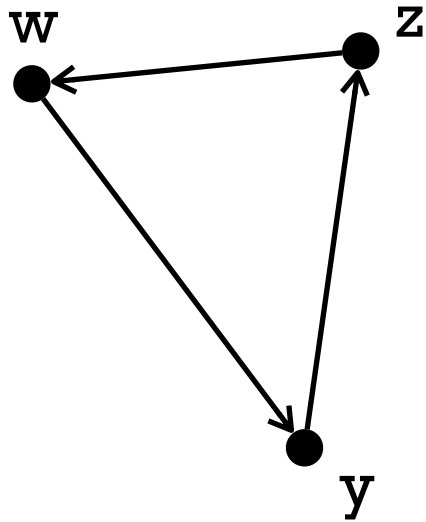


$z(c)$

$w(z(c))$



# (a) Acyclicity

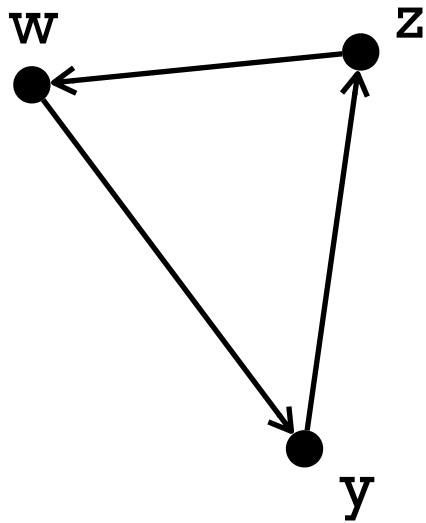


$z(c)$

$y(w(z(c)))$

$w(z(c))$

# (a) Acyclicity



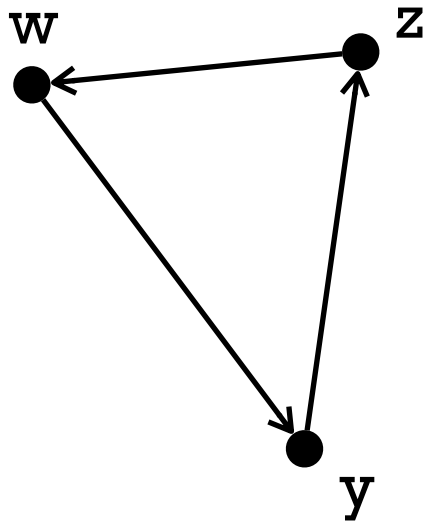
$z(c)$

$y(w(z(c)))$

$w(z(c))$

$z(y(w(z(c))))$

# (a) Acyclicity



$z(c)$

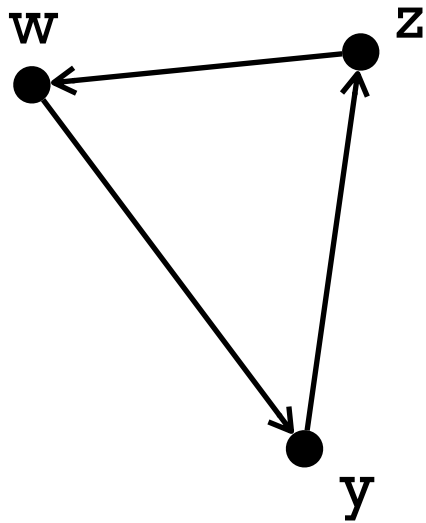
$y(w(z(c)))$

$w(z(c))$

$w(z(y(w(z(c))))))$

$z(y(w(z(c))))$

# (a) Acyclicity



$z(c)$

$y(w(z(c)))$

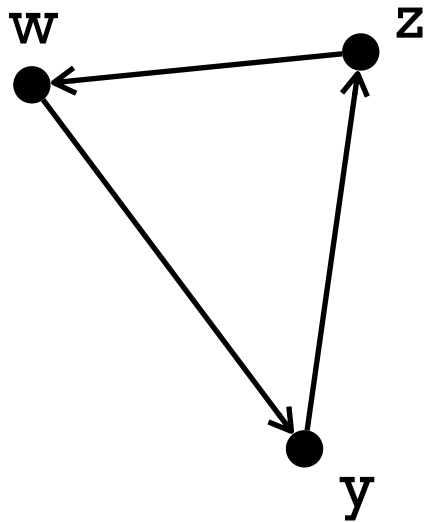
$w(z(c))$

$w(z(y(w(z(c))))))$

$z(y(w(z(c))))$

.....

# (a) Acyclicity



*Remark.* If the existential dependency graph of a given set of rules is acyclic, then the set of terms introduced during the computation of the chase is finite.

## (f) Arity at Most 1

$\text{Film}(x) \rightarrow \exists y . \text{IsFilmDirectedBy}(x, y) \wedge \text{Director}(y)$

$A(x) \wedge B(x, w) \wedge C(x, z) \rightarrow \exists z . R(x, w, z)$

## (f) Arity at Most 1

$\text{Film}(x) \rightarrow \exists y . \text{IsFilmDirectedBy}(x, y) \wedge \text{Director}(y)$

$\text{Film}(x) \rightarrow \text{IsFilmDirectedBy}(x, y(x)) \wedge \text{Director}(y(x))$

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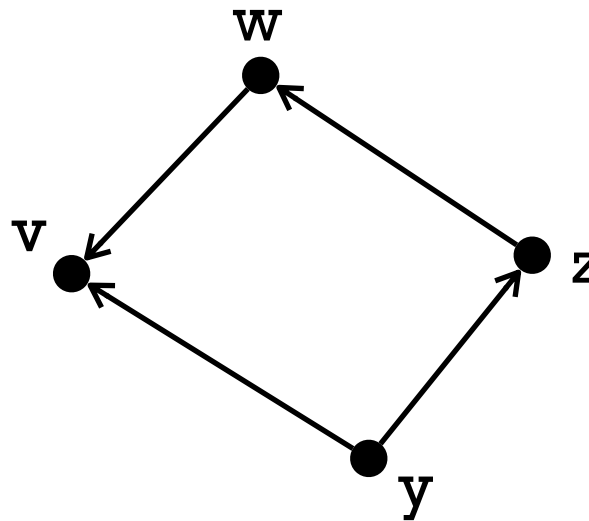
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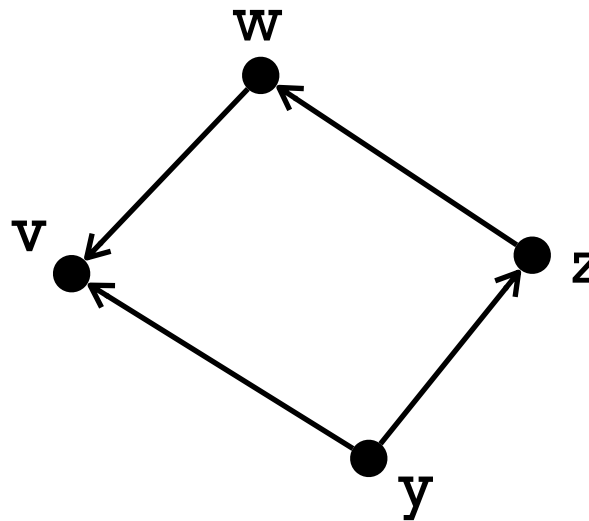
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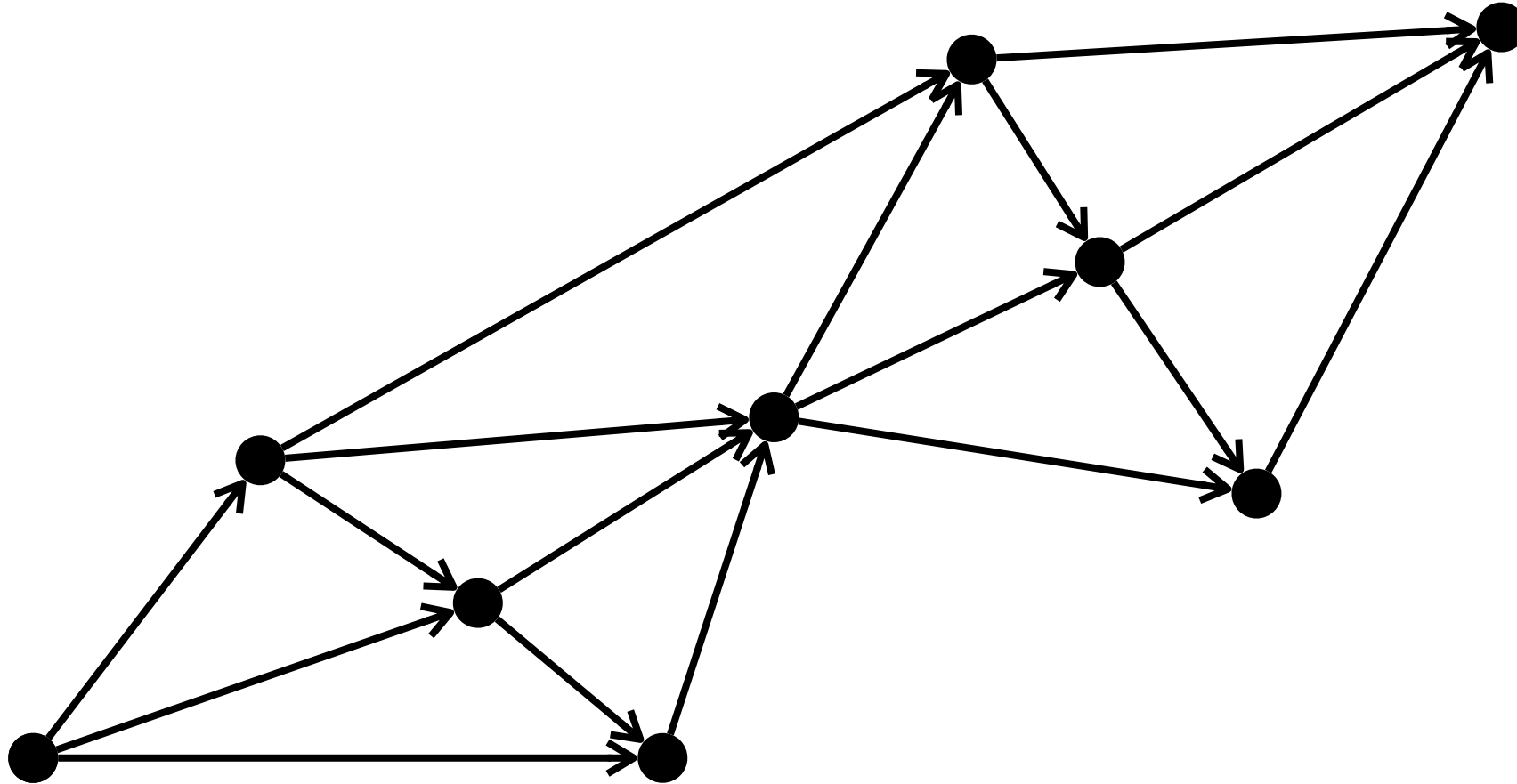
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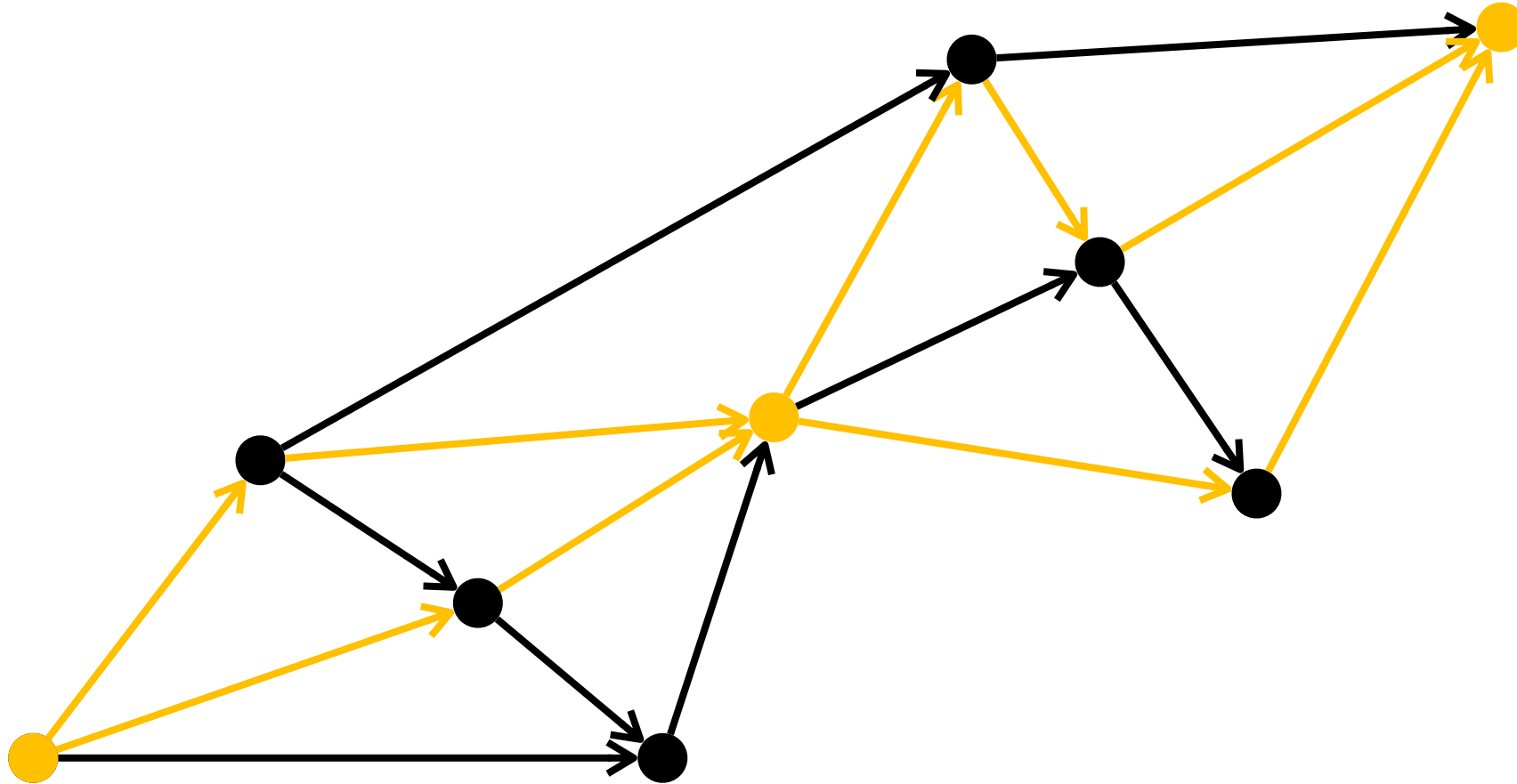
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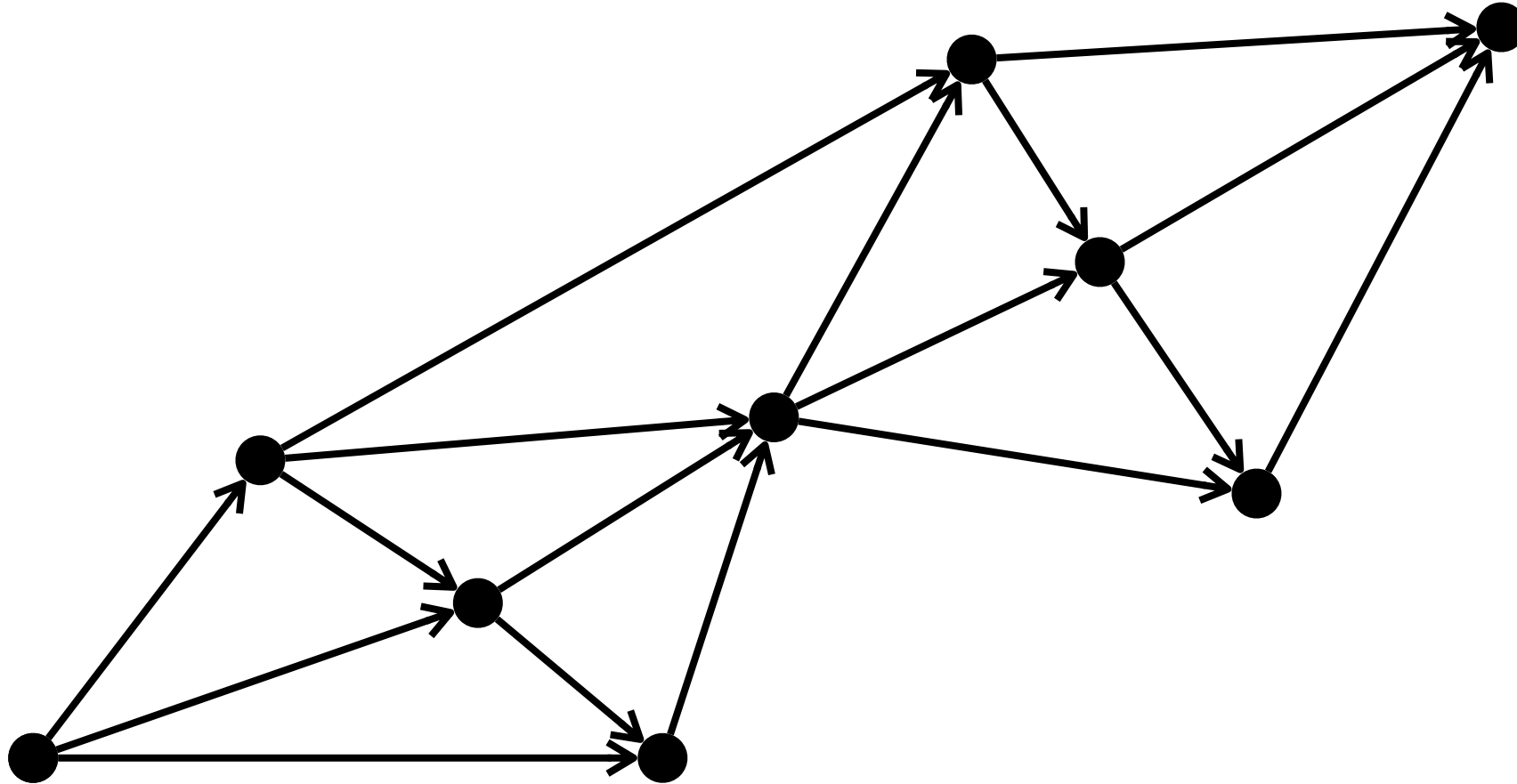
# Braids



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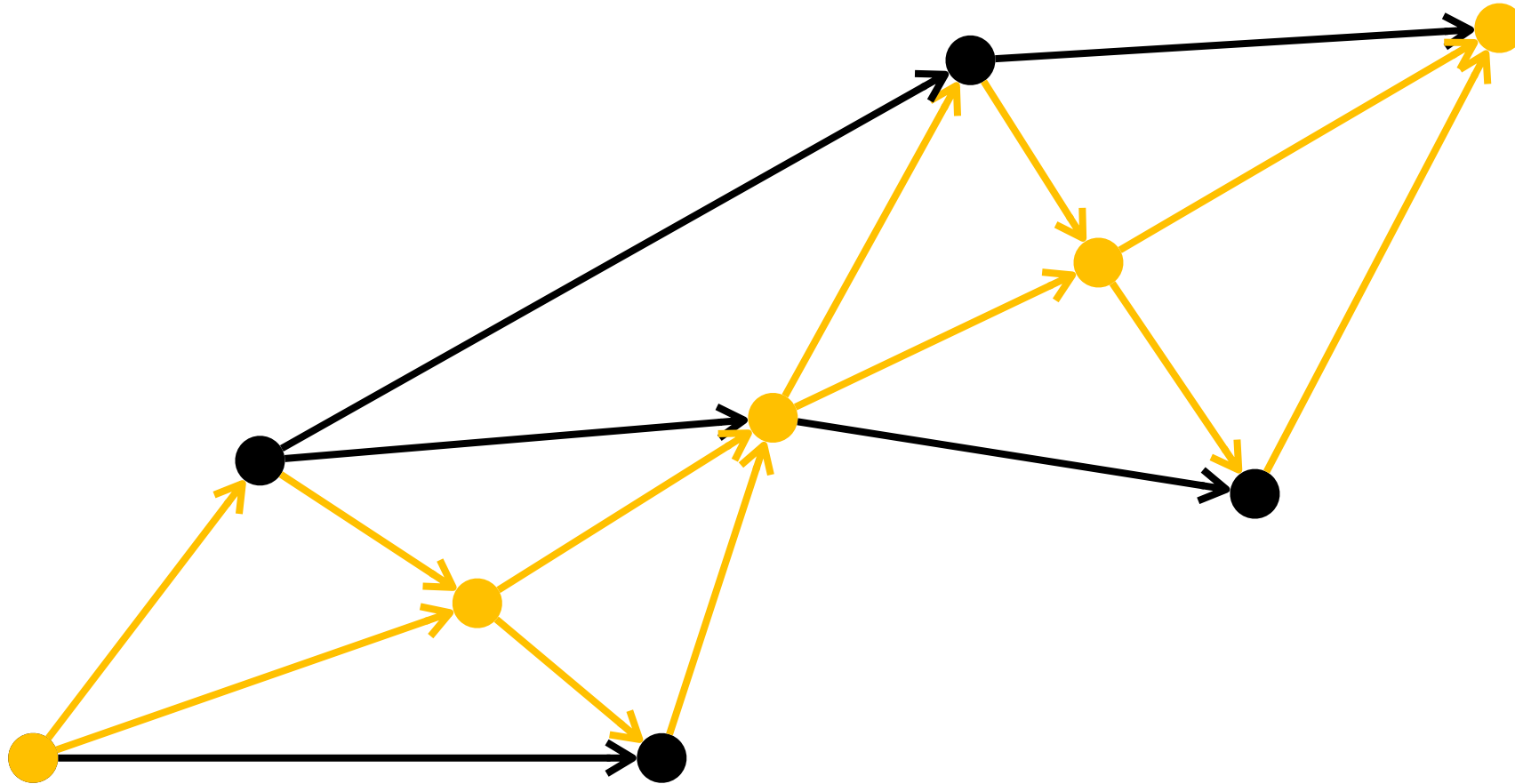


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## Caveats.

1. Fixed query size.
2. Horn rule set.

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# Evaluation

# SRI Axioms

$$\begin{aligned}A_1 \sqcap \dots \sqcap A_n \sqsubseteq B &\mapsto A_1(x) \wedge \dots \wedge A_n(x) \rightarrow B(x) \\A \sqsubseteq B_1 \sqcup \dots \sqcup B_n &\mapsto A(x) \rightarrow B_1(x) \vee \dots \vee B_n(x) \\A \sqsubseteq \forall R. B &\mapsto A(y) \wedge R(x, y) \rightarrow B(x) \\A \sqsubseteq \exists R. B &\mapsto A(x) \rightarrow \exists y. R(x, y) \wedge B(y) \\R \sqsubseteq S &\mapsto R(x, y) \rightarrow S(x, y) \\R \circ S \sqsubseteq V &\mapsto R(x, y) \wedge S(y, z) \rightarrow S(x, z) \\R_1 \sqcap \dots \sqcap R_n \sqsubseteq S &\mapsto R_1(x, y) \wedge \dots \wedge R_n(x, y) \rightarrow S(x, y) \\A(a) &\mapsto A(a) \\R(a, b) &\mapsto R(a, b)\end{aligned}$$

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- 1. Every rule in a SRI ontology has at most 3 variables.*
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*Corollary.* To guarantee that tractable reasoning over a SRI ontology is possible we only need to verify the following:

- 1. Acyclicity.*
- 2. Braid length in the dependency graph is bounded.*

# Evaluation Results

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## Acyclicity

	MOWL Corpus	Oxford Ontology Repo
Ontologies	1576	225
Acyclic	974 (61.8%)	170 (75.6%)



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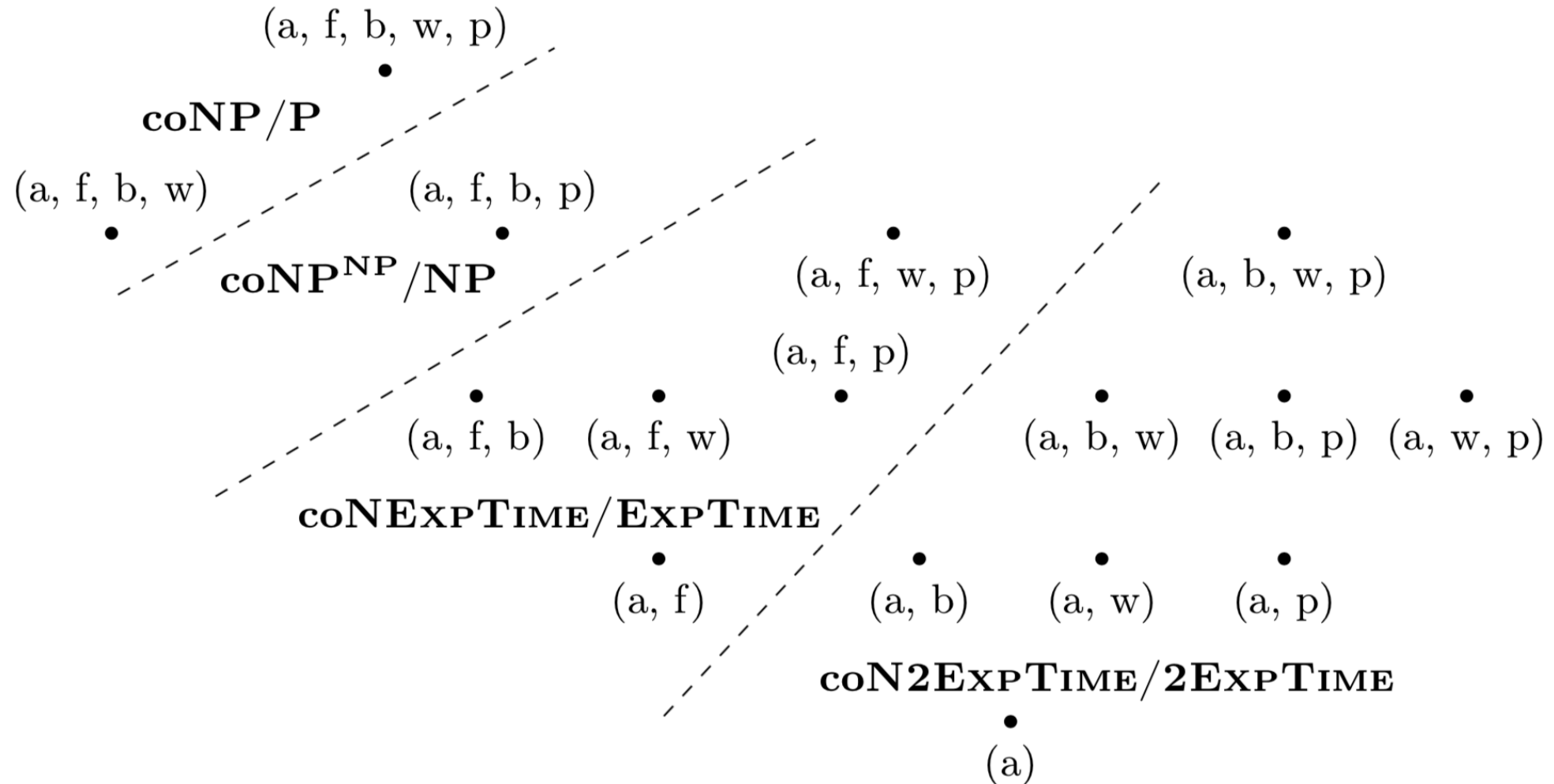
## Braid Length

### MOWL Corpus + Oxford

(max. length of a braid)	1	2	3	4	5	6	11	22	23	25	Total
(count)	851	153	56	61	11	1	1	2	7	1	1144
(%)	74	88	93	98	99	99	99	99.1	99.3	99.9	100

# Conclusions

# More Results!



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Efficient CQ Answering over a large subset of OWL 2 real-world ontologies is possible!

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Future work:

- Implementation of a chase based reasoner for these ontologies
- Refine conditions

# Tractable Query Answering for Expressive Ontologies and Rules\*

David Carral, Irina Dragoste, Markus Krötzsch  
Knowledge-Based Systems Group at



\*Published at ISWC 2017