

Logic on MARS: Ontologies for Generalised Property Graphs

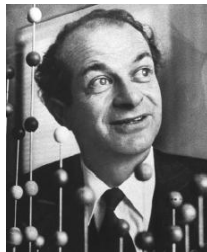
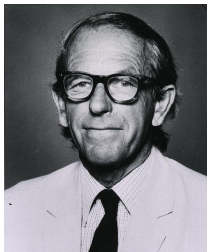
Maximilian Marx Markus Krötzsch Veronika Thost

TU Dresden

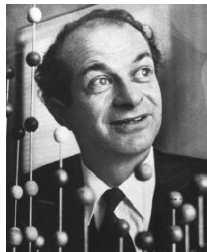
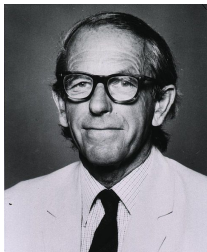
IJCAI 2017

Full paper: <https://iccl.inf.tu-dresden.de/web/MARS/en>

What do these people have in common?



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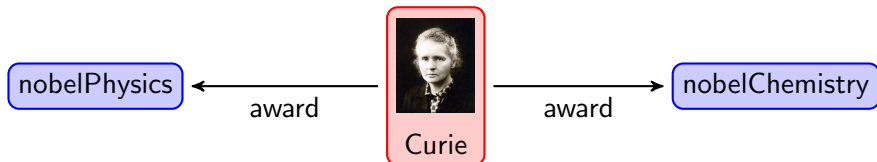


People that have won two Nobel prizes

Laureates in Knowledge Graphs



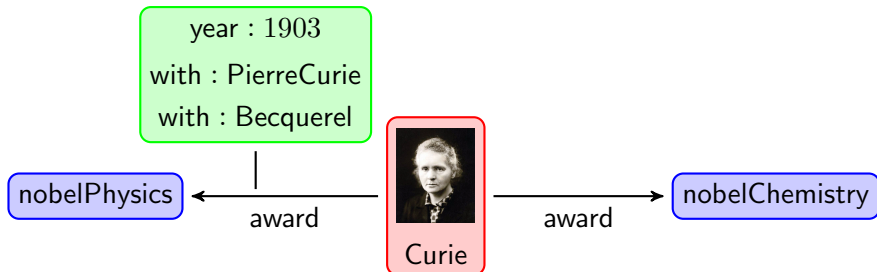
Wikidata: a free and open Knowledge Graph



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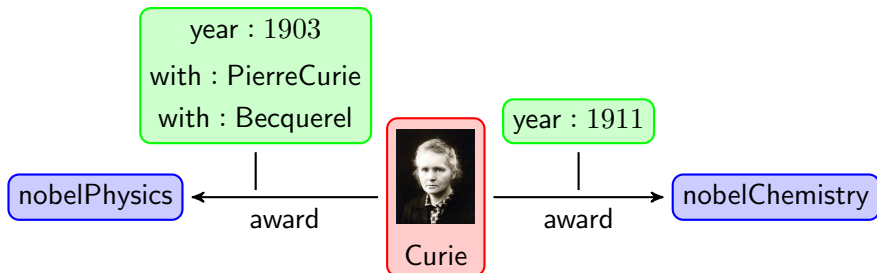
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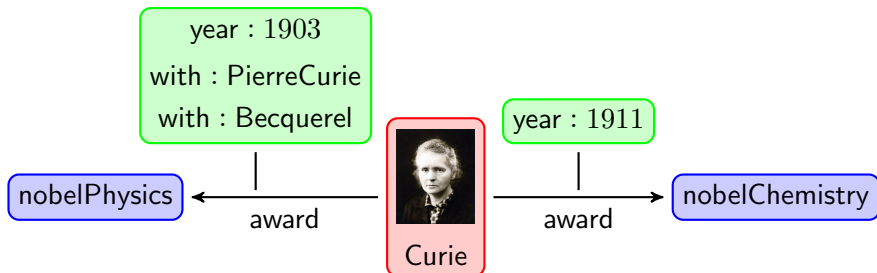
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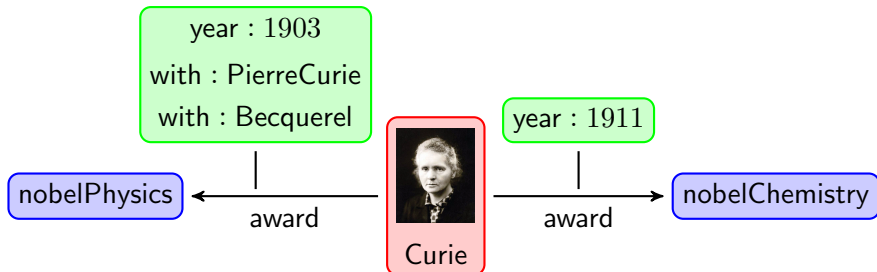


fact notation: `award(Curie, nobelChemistry)@{year : 1911}`

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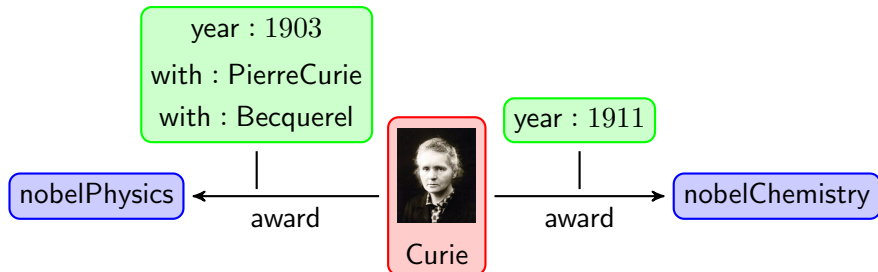
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Knowledge Graph: labeled graph; edges carry **annotation sets** (finite sets of **attribute–value** pairs)

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MARS: multi-attributed relational structure (annotated hypergraph)

Annotation-aware reasoning

Example: for sole “award” winners, infer “laureate,” i.e.,

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$$\forall x, y. \forall S. \text{award}(x, y)@S \wedge ([] \setminus [\text{with} : +])(S) \rightarrow \text{laureate}(y, x)@S$$

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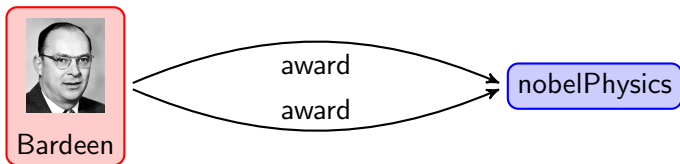
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Specifier $([] \setminus [\text{with} : +])$: any annotation set without attribute “with”

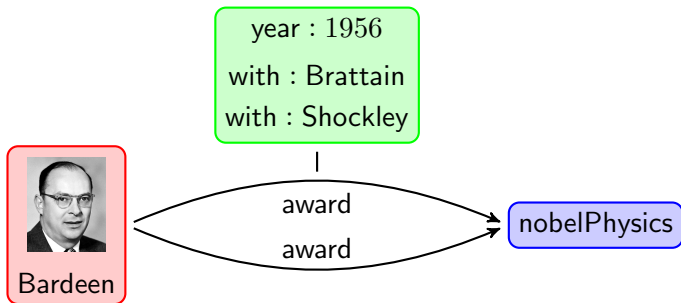
✓ {year : 1911}

✗ {year : 1903, with : PierreCurie, with : Becquerel}

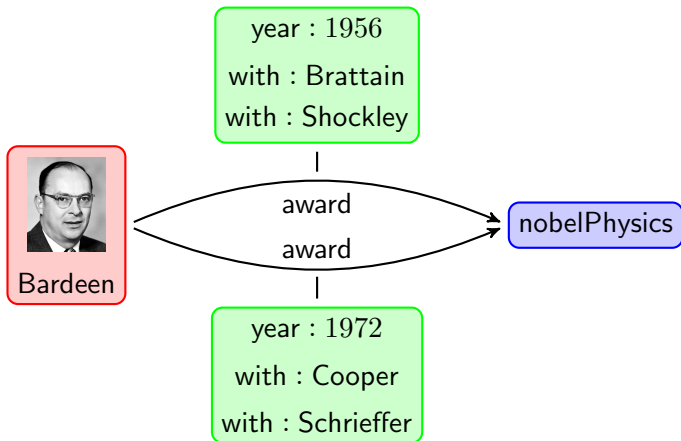
What else do we need?



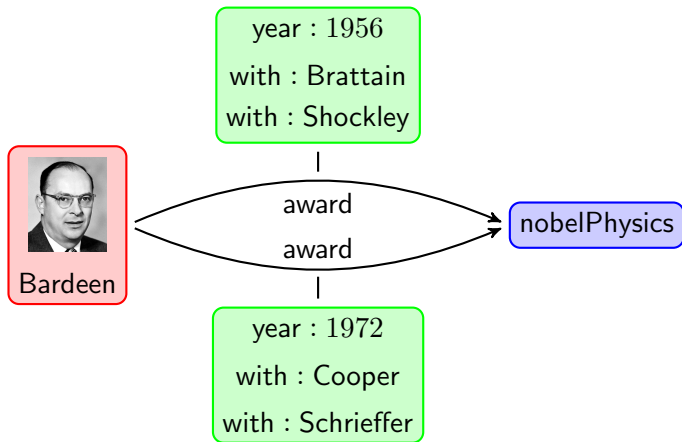
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Goal: derive “award” for co-laureates \rightsquigarrow copying annotation sets is not enough, we need to compute a new annotation set

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- ▶ derived knowledge may depend on annotation sets

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Theorem

MAPL has the same expressivity as weak second-order logic.

↪ Entailment for MAPL theories is not semi-decidable.

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Idea: Encode arbitrary arity predicates in annotation sets

A decidable fragment

- ▶ **Specifiers** express constraints on annotation sets

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\rightsquigarrow **Rule**

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with function definition **CoLaureate**(U, v, w):

$\Rightarrow \text{insert}(\text{with} : v)$

$\llbracket \text{with} : o \rrbracket(U), o \neq w \Rightarrow \text{insert}(\text{with} : o)$

$\llbracket \text{year} : d \rrbracket(U) \Rightarrow \text{insert}(\text{year} : d)$

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MARPL entailment is PTIME -complete for data complexity if the size of annotation sets is bounded.

Conclusion & Outlook

Summary:

- MAPL general, second-order based framework for attributed logics; not semi-decidable
- MARPL decidable, rule-shaped fragment; EXPTIME-complete for data & combined complexity

Future Work:

- ▶ Create attributed ontologies, e.g., for Wikidata
- ▶ Implement a MARPL reasoner
- ▶ Identify more expressive decidable fragments of MAPL
- ▶ Study attributed versions of other KR formalisms
- ▶ Classify data complexities, identify tractable fragments